## CHAPTER 21

## PRACTICE SET

## Questions

Q21-1. In multicasting, the sender host sends only one copy of the message, but it is multiplied at the routers if needed; all multiplied copies have the same destination address. In multiple-unicasting, the sender host sends one copy for each destination; each copy has its own destination address.

Q21-3. The multicast address block is 224.0.0.0/4. In other words, a multicast address is between 224.0.0.0 and 239.255.255.255. Based on this criteria we have
a. A multicast
b. A multicast
c. Not a multicast

Q21-5. If a host is a member of $N$ multicast group, it will have $N$ multicast addresses.

Q21-7.
a. In unicast communication, the destination is only one of the leaves of the tree in each transmission.
b. In multicast communication, the destination may be one or more leaves of the tree in each transmission.

Q21-9. DVMRP allows a router to create the shortest path-tree whenever it receives a multicast packet (on demand).
a. The number of shortest-path trees in DVMRP that use the source-based approach is huge.
b. If each router created all of the required multicast shortest-path trees, it would be a huge overhead.

Q21-11. RPB creates a broadcast shortest path tree with the source as the root and networks as the leaves. In other words, it tries a shortest path from the source to every network assuming that all networks are interested to receive that particular group message.

Q21-13. Each router using DVMRP creates the shortest-path three in three steps:
a. In the first step, the router uses the RPF algorithm to keep only packets that have arrived from the source using the shortest-path three. In other words, the first part of the tree is made using the RPF algorithm.
b. In the second step, the router uses the RPB algorithm to create a broadcast tree.
c. In the third step, the router use the RPM algorithm to change the broadcast tree created in the second step to a multicast tree.

Q21-15. Every multicast routing algorithm needs to somehow use a unicast protocol in its operation. For example, DVMRP needs to use RIP and MOSPF needs to use OSPF. Although PIM also needs to use a unicast protocol, the protocol can be either RIP or OSPF.

Q21-17. In PIM-DM, it is assumed that most networks have a loyal member in each group, so it does not matter if the first packet reaches all networks. In PIMSM, it is assumed that a few networks has a loyal member in each group, so broadcasting is wasting the bandwidth.

## Problems

P21-1. We define two properties for a block in classless addressing. The first address needs to divide the number of addresses in the block. The number of addresses in the block should also be a power of 2 . Let us check these properties for these two blocks (two applets for chapter 4 can easily find the values).
a. In the first block, the value of the first address is $3,758,096,895$ and the number of addresses in the block is 65,024 . None of the conditions meet. The block is intended to be split in the future.
b. In the second block, the value of the first address is $3,758,292,992$ and the number of addresses in the block is $134,021,120$. None of the conditions meet. The block is intended to be split in the future.
c. In the third block, the value of the first address is $3,925,868,544$ and the number of addresses in the block is $83,886,080$. None of the conditions meet. The block is intended to be split in the future.

P21-3. We use binary notation to do transformation:

|  |  | 11011111 | 00011000 | 00111100 | 00001001 |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  | $\downarrow$ | $\downarrow$ | $\downarrow$ |  |
| 00000001 | $\underline{00000000}$ | $\underline{01011110}$ | $\underline{00011000}$ | 00111100 | 00001001 |

The resulting address in hexadecimal is: 01:00:5E:18:3C:09. Using the same method, we get the same result.

P21-5. There is no need for a report message to travel outside of its own network because its only purpose is to inform the next router in the spanning tree of group membership. There is no need for a query message to travel outside of the local network because its only purpose is to poll the local network for membership in any groups.

P21-7. Please correct the errors in this problem using the errata. Since router B tells that the source is 4 hops away, router $B$ needs to be selected as the designated router.

P21-9. The following show the shortest path from the source and the shortest path tree seen from router R.

a. Shortest path tree seen by router S

b. Shortest path tree seen by router R

