PRACTICE SET

Questions

- Q22-1. We can say (1) larger address space, (2) better header format, (3) new options, (4) allowance for extension, (5) support for resource allocation, and (6) support for more security.
- Q22-3. A compatible address is an address of 96 bits of 0s followed by 32 bits of an IPv4 address. A mapped address is an address of 80 bits of 0s followed by 16 bits of 1s and followed by 32 bits of an IPv4 address. A compatible address is used when a computer using IPv6 wants to send a packet to another computer using IPv6. A mapped address is used when a computer using IPv6 wants to send a packet to a computer still using IPv4.
- Q22-5. The IP header is included because it contains the IP address of the original source. The first 8 bytes of the data are included because they contain the first section of the TCP or UDP header which contains information about the port numbers (TCP and UDP) and sequence number (TCP). This information allows the source to direct the ICMP message to the correct application.
- **Q22-7.** The prefix length is 3, which means the size of the block is

 $N = 2^{128-3} = 2^{125}$ A very large block

Q22-9. The prefix length is 7, which means the size of the block is

 $N = 2^{128-7} = 2^{121}$ A very large block

- **Q22-11.** *Autoconfiguration* allows the DHCP to automatically allocate an IP address to a host when the host joins the network.
- Q22-13. The *next header* field is responsible for multiplexing and demultiplexing in IPv6. It is similar to the protocol field in IPv4.

Q22-15.

- **a.** The *neighbor-solicitation* message in ICMPv6 replaces the ARP request message in version 4.
- **b.** The *neighbor-advertisement* message in ICMPv6 replaces the ARP response message in version 4.
- **Q22-17.** We need to encapsulate IPv6 packets in IPv4 packets in the *tunneling* strategy.

Problems

P22-1. The following table shows the comparison:

| Field | IPv4 | IPv6 |
|----------------------------|------|------|
| VER | Ń | Ń |
| HLEN | Ń | |
| Service (or traffic class) | Ń | Ń |
| Flow label | | Ń |
| Total length | Ń | |
| Payload length | | Ń |
| Identification | Ń | |
| Flags | Ń | |
| Flag offset | Ń | |
| TTL (or hop limit) | Ń | Ń |
| Protocol | Ń | |
| Checksum | Ń | |
| Source Address | Ń | Ń |
| Destination Address | Ń | Ń |

P22-3.

- a. 0:FFFF:FFFF::
- b. 1234:2346:3456::FFFF
- c. 0:1:::FFFF:1200:1000
- d. :: FFFF: FFFF: 24.123.12.6

P22-5.

- **P22-7.** Both subnets keep the given *global routing prefix*. Each subnet adds a 16-bit subnet identifier. We assume the subnet identifiers start from $(0001)_{16}$, but they can also start from $(0000)_{16}$.
 - **a.** The first subnet block is 2000:1234:1423:0001/64.
 - b. The second subnet block is 2000:1234:1423:0002/64.
- **P22-9.** We change the seventh bits from 0 to 1 (F5 is changed to F7) and insert four extra hexadecimal digits (FF-FE) after the sixth digits:

F5 A9 -23 -12 -7A -B2 F7 A9 23 FF FE 12 **7**A **B2**

The resulting preface address in IPv6 is F7A9:23FF:FE12:7AB2.

- P22-11. The address is ::129.6.12.34/128.
- **P22-13.** The address is ::01/128.
- **P22-15.** Here we mean unique local unicast address (see the errata). The prefix in this case is FC00/7. After adding the suffix (0::123/48), we can get the *unique local unique address* as FC::123/48.
- **P22-17.** See the following figure.

| 6 | 0 | | 3245 | | | |
|---------------------------|------------|---|------|----|--|--|
| | 0 | | 0 | 15 | | |
| 581E:14562314:ABCD::1211 | | | | | | |
| 581E:1456:2314:ABCD::2211 | | | | | | |
| | 6 | 8 | 194 | 4 | | |
| 128020 | | | | | | |
| | 60333 25 | | | | | |
| 345 | | | | | | |
| 7 | | | | | | |
| 5 | Reserved | 0 | 8192 | 2 | | |
| | Checksum 0 | | | | | |
| 128000 bytes of data | | | | | | |

P22-19. The error-reporting messages in ICMPv4 and ICMPv6 are similar except that some messages have been totally deleted in version 6 or has been inserted in other categories. The following table shows a comparison.

| Message | v4 | vб | Explanation |
|-------------------------|--------------|----|----------------------------------------|
| Destination unreachable | \checkmark | | |
| Source quench | \checkmark | | Deleted from ICMPv6; rarely used |
| Time exceeded | \checkmark | | |
| Parameter problem | \checkmark | | |
| Redirection | \checkmark | | Moved to neighbor-discovery category |
| Packet too big | | | Added in v6 to prevent big-packet size |

P22-21. The neighbor-discovery category is new in ICMPv6. In ICMPv4, one of these messages belonged to another category; the duties of others were covered by other protocols, The following table shows a comparison.

| Message | v4 | vб | Explanation |
|------------------------|----|--------------|-----------------------------------------|
| Router solicitation | | \checkmark | Version 4 uses DHCP for this purpose |
| Router advertisement | | \checkmark | Version 4 uses DHCP for this purpose |
| Neighbor solicitation | | \checkmark | Version 4 uses ARP request packet |
| Neighbor advertisement | | \checkmark | Version 4 uses ARP reply packet |
| Redirection | | \checkmark | Included in error-reporting group in v4 |

P22-23. Group membership messages are new in Version 6. Version 4 is using IGMP protocol for this purpose. The following table shows a comparison.

| Message | v4 | vб | Explanation |
|-------------------|----|--------------|-------------------------------------|
| Membership query | | \checkmark | It was part of IGMP protocol in v4. |
| Membership report | | \checkmark | It was part of IGMP protocol in v4. |